Reconstructing the history of a WD40 beta-propeller tandem repeat using a phylogenetically informed algorithm

Philippe Lavoie-Mongrain¹, Mahdi Belcaid², Aïda Ouangraoua^{3,4}, Anne Bergeron¹ and Guylaine Poisson⁵

- ¹ LaCIM, Université du Québec à Montréal, Canada
- ² HIMB, University of Hawaii at Manoa, Honolulu, USA
- ³ INRIA, LIFL, Université Lille 1, France
- ⁴ Computer Science, Université de Sherbrooke, Canada
- ⁵ ICS, University of Hawaii at Manoa, Honolulu, USA

ABSTRACT

Tandem repeat sequences have been found in great numbers in proteins that are conserved in a wide range of living species. In order to reconstruct the evolutionary history of such sequences, it is necessary to develop algorithms and methods that can work with highly divergent motifs. Here we propose a reconstruction algorithm that uses, in parallel, ortholog tandem repeat sequences from *n* species whose phylogeny is known, allowing it to distinguish mutations that occurred before and after the first speciation. At each step of the reconstruction, both the boundaries and the length of the duplicated segment are recalculated, making the approach suitable for sequences for which the fixed boundary hypothesis may not hold. We use this algorithm to reconstruct a 4-bladed ancestor of the 7-bladed WD40 beta-propeller, using orthologs of the GNB1 human protein in plants, yeasts, nematodes, insects and fishes. The results obtained for the WD40 repeats are very encouraging, as the noise in the duplication reconstruction is significantly reduced.

Subjects Bioinformatics, Computational Biology **Keywords** Ancestor reconstruction, Tandem repeats, Phylogeny, Beta propellers

INTRODUCTION

Among the many evolutionary events that shape genes, tandem repeats play a major role as building blocks of protein domains that have structural functions in a cell, and some very useful architectures have been conserved through a wide range of species (*Schaper*, *Gascuel & Anisimova*, 2014). One striking example is the 7-bladed *beta-propeller* of the human GNB1 protein which has orthologs in yeasts and plants: the repetitive nature of this structure is quite obvious in JSmol (*Hanson et al.*, 2013) views such as the one shown in Fig. 1. This repetition is mirrored by the underlying coding sequence that contains nearly 7 repetitions of the WD40 motif, which are short segments of approximately 40 amino acids, featuring a highly conserved **W-D** dipeptide.

Variation in the number of these building blocks may have an impact on a wide array of geometric parameters such as the diameter of cross-membrane channels, or the area and length of binding surfaces (*Andrade, Perez-Iratxeta & Ponting, 2001*). For example,

Submitted 13 February 2015 Accepted 12 May 2015 Published 27 May 2015

Corresponding author Anne Bergeron, bergeron.anne@uqam.ca

Academic editor Claus Wilke

Additional Information and Declarations can be found on page 15

DOI 10.7717/peerj-cs.6

© Copyright 2015 Lavoie-Mongrain et al.

Distributed under Creative Commons CC-BY 4.0

OPEN ACCESS

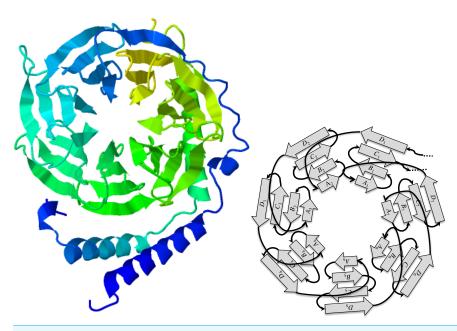


Figure 1 The GNB1 beta-propeller. In this JSmol (*Hanson et al., 2013*) ribbon diagram view of a GNB1 beta-propeller ortholog (PDB:1TBG; GenPept:P62871), the blades are displayed as seven triangles organized in a circle, like a propeller. Each blade is composed of four strands, labeled *A*, *B*, *C*, and *D* starting from the inner one. Reading the protein sequence from N-terminus to C-terminus, the first strand is the *D* strand of the upper right blade (D_1), followed by the *A*, *B*, *C* strands of the upper left blade. The *A*, *B*, *C* strands of the upper right blade (A_1 , B_1 and C_1) are the last strands of the protein sequence. This implies that, sequence-wise, blades are shifted with respect to the corresponding 3D units.

the number of **W-D** repeats in a domain typically ranges from 4 to 16 blades (*Li* & *Roberts*, 2001), making up a complex and interesting class of domains whose evolutionary relationships remain unresolved.

The goal of this paper is to propose algorithms and heuristics to reconstruct the evolutionary history of these ancient domains. The main challenges are the lack of sequence similarity between repeats, and the recognition that repeats may not evolve according to a tree-like model—which is a crucial hypothesis in most current approaches.

When form is more important than content, similarity between repeated components can degrade sufficiently to fool computational tools designed to predict them, some of which had to wait for their 3D structural resolution in order to be detected (*Kajava*, 2012). Reconstructing evolutionary histories of such repeats is particularly difficult, since it relies even more on similarity detection than the prediction problem does.

A second problem in the reconstruction of evolutionary histories is whether the boundaries of the repeated segment remained fixed throughout evolution, or varied from one duplication to the other. The fixed boundaries hypothesis allows for the use of alignments and/or phylogenetic trees that can give global solutions (*Gascuel, Bertrand* & *Elemento, 2005; Sammeth & Stoye, 2006; Tremblay Savard, Bertrand & El-Mabrouk, 2011*), but the applications of these techniques are restricted to tandem repeats constrained by biological features, such as the presence of introns separating the repeated segments (*Zhuo et al., 2007*).

To circumvent these problems, and since ancient domains are likely to have orthologs in several species, we use sets of orthologs to reconstruct, in parallel, their common evolutionary history. The backbone of the method (*Benson & Dong, 1999*) is the only known algorithm that does not rely on the fixed boundary hypothesis (*Rivals, 2004*). The Benson and Dong algorithm is parsimony-based, and we extend it using the same principle, given the computationally intensive nature of our approach.

We illustrate our approach using a set of six GNB1 ortholog **W-D** domains in mammals, plants, yeasts, nematodes, insects and fishes: these were shown to share their evolutionary history in at least two independent studies (*Chaudhuri, Soding & Lupas, 2008; Schaper, Gascuel & Anisimova, 2014*). We give evidence that the fixed boundary hypothesis might not apply, and we reconstruct a 4-bladed ancestor of the 7-bladed **W-D** GNB1 beta-propeller that matches living fossils such as slime molds.

METHODS

Preliminaries

A *tandem repeat* is a sequence composed of the repetition of approximate copies of a *motif*, where the first or the last repetition may be truncated. For example, the tandem repeat:

CC GTTAC GTTACC GTAGC

contains approximate repetitions, separated by spaces, of the motif GTTACC.

When the motif and the copies have the same length, we call this length a *period*. With insertions and deletions in the copies, it is always possible to transform a tandem repeat such that all the complete copies have the same length. For example, the above sequence can be rewritten, by inserting two gaps, as: CC GTTA-C GTTACC GT-AGC. Once all copies have the same length, the trivial multiple alignment of the successive copies is called a *self-alignment*.

The problem of self-aligning a tandem repeat sequence while optimizing any interesting score—sum-of-pairwise-distances, star and tree alignments, etc.—is NP-hard (*Elias*, 2006), and is at the heart of every reconstruction strategy (*Rivals*, 2004). However, in the case of tandem repeats that code for protein domains whose structure is resolved, such as in Fig. 1, structure-guided alignments can be used to infer the corresponding self-alignments (*Chaudhuri, Soding & Lupas, 2008*).

The *history* of a tandem repeat is a description of how it went from one copy of the motif in the ancestor to the many copies in the observed sequence. This can be formalized using various models of evolution. All these models include mutations, insertions and deletions of nucleotides that affect a single copy of the motif, and *duplications* that replace a segment of *k* copies of the motif by two consecutive segments both identical to the initial segment (*Benson & Dong, 1999*). More elaborate models will consider operations that transform a segment into two or more identical segments (*Gascuel, Bertrand & Elemento, 2005*), deletion of copies (*Sammeth & Stoye, 2006*), or even inversions and speciations (*Tremblay Savard, Bertrand & El-Mabrouk, 2011*). The main hurdle in the reconstruction problem is the fact that more than a billion years of evolution can

significantly blur the similarity of the sequences, thus rendering quantitative comparison tools meaningless.

Algorithmic approaches vary according to the evolutionary model, and more fundamentally according to assumptions on the behavior of duplications, linked to the concept of *parsing*. Given a period *p*, a *parsing* of a tandem repeat is the determination of the *boundaries* of the copies of a repeated motif of length *p* (*Matroud et al., 2012*). For example, here are two different parsings of the example sequence:

CC | GTAA-C | GTTACC | GT-AGC CCGTA | A-CGTT | ACCGT- | ACC

In the *fixed* boundaries model, the boundaries of all duplications are boundaries of the same parsing; in the *dynamic* boundaries model, parsing may vary from one duplication to the next. With the fixed boundaries hypothesis the reconstruction problem can be reformulated as the construction of a *duplication* tree, opening up the vast repertoire of phylogenetic tools. In contrast, the more general dynamic boundaries hypothesis implies that exhaustive searches must be repeated at each step (*Benson & Dong, 1999*), and the evolutionary history cannot be represented by a tree (*Jaitly et al., 2002*).

Architecture of the GNB1 beta-propeller

The *guanine nucleotide-binding protein* (GNB1) is a protein characterized by a tandem repeat of seven copies of the WD40 protein motif (PFAM00400). These fold in a three-dimensional cylindrical structure, whose top view, as seen in Fig. 1, looks like a propeller with seven blades. Each blade is a beta sheet consisting of four strands, denoted *A*, *B*, *C*, and *D*, from the centre of the propeller to its circumference. However, the strands appear in the tandem repeat in the following order:

$D_1A_2B_2C_2D_2A_3B_3C_3D_3A_4B_4C_4D_4A_5B_5C_5D_5A_6B_6C_6D_6A_7B_7C_7D_7A_1B_1C_1$

Thus, there are at least two competing parsings of the tandem repeat: one that preserves the integrity of the 3D blades, with motif *ABCD*, and one that maximizes the number of complete copies, with motif *DABC*. In view of this particularity, we reject the fixed boundaries hypotheses in favor of dynamic boundaries, leaving open the possibility of eventually detecting fixed boundaries.

The WD40 protein motif, which actually codes for the components of a blade, is quite common (*Stirnimann et al., 2010*), and is responsible for a variety of propellers ranging from 4 to 8 blades, whose ancestry can be traced back to both recent rounds of duplications, and very ancient ones (*Chaudhuri, Soding & Lupas, 2008*). The GNB1 propeller falls in the latter category, and has been found in all sequenced eukaryotes species.

A GNB1 ortholog tandem repeats dataset

We selected a set of GNB1 orthologs using the Ensembl (*Flicek, Amode & Barrell, 2014*) database, and BLAST (*Altschul et al., 1990*) in the case of Arabidopsis. Table 1 gives

Code	Species	Accession	Positions
HUM	Homo sapiens	NP_002065	48-340
DAN	Danio rerio	NP_997774	48-340
DRO	Drosophila melanogaster	NP_996462	48-340
CEL	Caenorhabditis elegans	AAK55963	48-340
NEU	Neurospora crassa	XP_956704	63–357
ARA	Arabidopsis thaliana	NP_001078491	53-364

 Table 1 Dataset. Sequences used to construct the dataset.

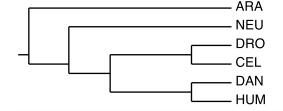


Figure 2 The phylogeny used by the algorithm. The species phylogeny from *Dunn et al. (2008)* used as input in the reconstruction algorithm.

the accession numbers and the positions of the WD40 tandem repeat in each sequence. Figure 2 shows the current consensus of the phylogeny of the six species (*Dunn et al., 2008*).

We used the structure-guided alignment of GNB1 given in *Chaudhuri, Soding & Lupas* (2008) to identify insertions and deletions of the human sequence within the WD40 motif. We then aligned the human sequence with *A. thaliana* and *N. crassa* sequences in order to transfer the insertions and deletions from the human, and to identify further deletions in *A. thaliana* and *N. crassa*. The alignment is shown in Fig. 3. The three remaining species inherited the same insertions and deletions as the human.

The resulting amino acid sequences have all the same length. These sequences are ortholog tandem repeats with 6.8 copies of the WD40 motif, as shown in the self alignment of the human sequence in Fig. 4. The underlying genomic sequences have a length of 864 bp; in the rest of this paper, they will be denoted by the names HUM, DAN, DRO, CEL, NEU, ARA, as in Table 1.

For comparison purposes, we computed the initial alignment with the standard tools Clustal Omega (*Sievers et al., 2011*) and Muscle (*Edgar, 2004*). Of the eleven or so mutations proposed by Chaudhuri et al., Clustal finds 7 of them, and Muscle only 2, we thus used the Clustal alignment. No classical aligner can predict, as Chaudhuri et al. do with structural knowledge, the two deletions that occur in all sequences. These deletions are necessary information for the computation of self-alignments.

Pinpointing the most recent duplication

Here we describe how to detect the most recent duplication in a tandem repeat, first by using species separately, then by processing the six species in parallel. The first model follows closely the *Benson & Dong* model (1999). The second is an enhanced model that

HUM	${\tt RRTLRGHLAKIYAMHWGTDSRLLVSASQDGKLIIWDSYTTNKVHAIPLRSSWVMTCAYAP$
NEU	KKTLKGHLAKIYAMHWSTDRRHLVSASQDGKLIIWDAYTTNKVHAIPLRSSWVMTCAYAP
ARA	CRTLQGHTGKVYSLDWTPERNRIVSASQDGRLIVWNALTSQKTHAIKLPCAWVMTCAFSP
	***** .******* * . ********************
HUM	SGNYVACGGLDNICSIYNLKTREGN vrvs RELAGHTGYLSCCRF-LD-DNQIVTSSG
NEU	SGNFVACGGLDNICSIYNLNSnRDGPtrvyRELSGHAGYLSCCRF-IN-DRSILTSSG
ARA	NGQSVACGGLDSVCSIFSLSS tad KDGT vpvs RMLTGHRGYVSCCQYVPNeDAHLITSSG
	·*: *****·:***:·*: ::* · * * *:** **:**::: * ::****
HUM	DTTCALWDIETGQQTTTFTGHTGDVMSLSLAPDT-RLFVSGACDASAKLWDVREG
NEU	DMTCMKWDIETGTKVVEFADHLGDVMSISLNPTN q NTFVSGACDSFAKLWDIRAG
ARA	DQTCILWDVTTGLKTSVFG <mark>gefqs</mark> GHTADVLSVSISGSN p NWFISGSCDSTARLWDTRAA
	* ** **: ** :. * * ***:*: *:**:**: *:*** * .
HUM	-MCRQTFTGHESDINAICFFPNGNAFATGSDDATCRLFDLRADQELMTYSHDNIIcgI
NEU	-KAVQTFAGHESDINAIQFFPDGHSFVTGSDDATCRLFDIRADRELNCYRSESILcgI
ARA	<mark>s</mark> RAVRTFHGHEGDVNTVKFFPDGYRFGTGSDDGTCRLYDIRTGHQLQVYQPH <mark>gd</mark> GEN <mark>gp</mark> V
	· :** ***·*:*: ***:* * ******:*:*:*:*: ::* * · · · ·
HUM	TSVSFSKSGRLLLAGYDD-FNCNVWDALKADRAGVLAGHDNRVSCLGVTDDGMAVA
NEU	TSVATSVSGRLLFAGYDD-FECKVWDLTRAEKVGSLVGHENRVSCLGVSNDGISLC
ARA	TSIAFSVSGRLLFAGYASn NTCYVWDTLLGEVVLD lglq QDSHRNRISCLGLSADGSALC
	:: * **:*** . * *** .:* **:***:: ** ::.
HUM	TGSWDSFLKIWN
NEU	TGSWDSLLKVWA
ARA	TGSWDSNLKIWA
	***** **:*

Figure 3 Alignment of three GNB1 ortholog sequences. The beta-propeller domain of the three GNB1 ortholog sequences from HUM, NEU and ARA were aligned using Clustal Omega (*Sievers et al., 2011*). The alignment was edited as shown: two blocks (green) were removed from all the sequences, according to the structural alignment of *Chaudhuri, Soding & Lupas (2008)* that suggested that these blocks do not belong to the structure; one gap (blue) was inserted in the HUM and NEU sequence, again according to *Chaudhuri, Soding & Lupas (2008)*; eight blocks (red) were removed from the NEU and ARA sequences corresponding to gaps in the HUM sequence. The remaining sequences, DAN, DRO and CEL, have exactly the same editions as the HUM sequence, since they align without gaps. The six resulting edited sequences have all the same length, 288 aa. The nucleotide sequences coding for these sequences were then extracted, to yield the main dataset of six sequences of length 864 bp.

uses phylogenetic trees of ortholog domains to distinguish mutations that occurred before all speciation events, from mutations that occurred after speciation.

Using a single species

After a duplication of a segment, the two copies are identical. If each copy accumulates mutations independently and uniformly, which is equivalent to a molecular clock hypothesis among the repeated segments, the most recent duplication should be indicated by the pair of segments that have the least differences. This leads to the following definition of a *cost* function, whose value is proportional to the number of mutations that separates two segments, and called the *contraction cost* (*Benson & Dong, 1999*). Let *S* be a nucleotide sequence, ℓ be a period length, and the *edit* distance given by $\delta(n_1, n_2) = 0$ if nucleotides n_1

HUM1	RRTLRGHLAKIYAMHWGTDSRLLVSASQDGKLIIWD
HUM2	SYTTNKVHAIPLRSSWVMTCAYAPSGNYVACGGLDNICSIYN
HUM3	LKTREGNRELAGHTGYLSCCRF-LDDNQIVTSSGDTTCALWD
HUM4	IETGQQTTTFT <mark>GH</mark> TGDVMSLSLAPDTRLFVSGAC <mark>D</mark> ASAKLWD
HUM5	VREGMCRQTFTGHESDINAICFFPNGNAFATGSDDATCRLFD
HUM6	LRADQELMTYSHDNIIITSVSFSKSGRLLLAGYD D FNCNVWD
HUM7	ALKADRAGVLAGHDNRVSCLGVTDDGMAVATGSWDSFLKIWN

Figure 4 Self-alignment of the human sequence. A self-alignment of HUM sequence shows a 42 amino acids long tandem repeat of 6.8 copies of the WD40 domain. The highlights of the WD40 repeat (*Stirnimann et al., 2010*) are illustrated by the high conservation of the **G-H** and **W-D** groups (in blue), and the conserved **D** residue (in red).

and n_2 are equal, otherwise $\delta(n_1, n_2) = 1$. We define the *contraction cost* of two consecutive segments of length ℓ at position *pos* of the sequence *S* as:

$$C(pos,\ell) = \frac{\sum_{i=pos}^{pos+\ell-1} \delta(S_i, S_{i+\ell})}{\ell}.$$
(1)

The normalization of the cost will allow the comparison of costs between different values of ℓ .

Given a motif of length p and a tandem repeat sequence S, the position of the most recent duplication of one or more copies of the motif should be a position in S that minimizes $C(pos, \ell)$, for all values of pos and $\ell \in \{k * p \mid k \ge 1\}$ for which $C(pos, \ell)$ is defined.

For the three GNB1 ortholog sequences HUM, NEU and ARA, we tested all possible periods and obtained minimal values for p = 126. Figure 5 shows the values of C(pos, 126) for each species. The prediction that no meaningful minimal values would emerge is confirmed both by the random nature of the HUM curve, and by the fact that, although these three sequences share the same duplication history, they do not agree on its most recent development.

Using multiple species

In this section, we use the same global approach as the preceding section, but we estimate the cost of a contraction based on the reconstruction of the ancestral tandem repeat of n extant tandem repeats who share the same duplication history, and whose phylogeny is known. We replace the distance $\delta(n_1, n_2)$ between nucleotides—whose value is 0 or 1—by a cost that ranges in the interval [0..1]. The following small example illustrates the method.

Consider the following three ortholog sequences, each with a tandem repeat of two repetitions of a motif of length 3.

	1	2	3	1'	2′	3'
Sp1	Т	A	G	Т	Т	Т
Sp2	С	A	A	A	Т	Т
Sp3	Т	С	С	A	Т	Т

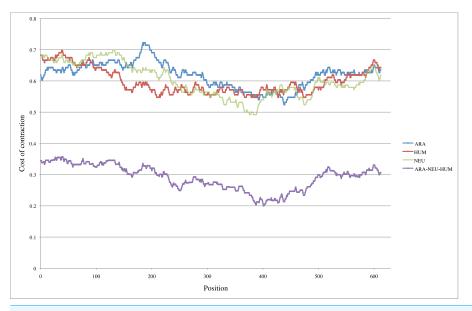


Figure 5 Pinpointing the most recent duplication. Top three curves: For each position *pos* in the three nucleotide sequences HUM, NEU and ARA, we computed the cost *C*(*pos*, 126) of contracting segments [*pos..pos* + 125] and [*pos* + 126..*pos* + 251]. The positions where minimal values are attained presumably indicate where the most recent duplication occurred. For the HUM sequence (red curve), positions corresponding to minimal values are scattered within the interval [209..477]. Although the intervals are smaller in the case of NEU and ARA (green and blue curves) those intervals do not even intersect. **Bottom curve:** This is the curve obtained by combining the data from HUM, NEU and ARA. The values of all costs are drastically reduced, with the absolute minimum reduced by more than half (0.492 for the NEU curve compared to 0.200 for the combined curve). The interval of positions where minimal values are attained is reduced to [387..403].

At positions 1 and 1', both species Sp2 and Sp3 have a mutation compared to Sp1. If the phylogeny is (Sp1, (Sp2, Sp3)), then all parsimonious reconstructions of the ancestral nucleotide give **T** for both positions 1 and 1' as in Tree 1 of Fig. 6. In the case of positions 2 and 2', all parsimonious reconstructions, such as Trees 2a and 2b of Fig. 6, give **A** for position 2, and **T** for position 2'.

For positions 3 and 3', some reconstructions have a mutation that precedes speciation and some do not. Figure 6 gives four possible reconstructions: Trees 3a, 3b and 3c have one mutation before speciation and Tree 3d does not.

Formally, given *n* ortholog tandem repeats, from *n* species whose phylogeny is given by tree \mathcal{T} , construct the tree \mathcal{D} whose root is a duplication node, with two copies of tree \mathcal{T} labeling its two children. The leaves of the left subtree are labeled by the nucleotides at position *pos* for each species, and the leaves of the right subtree are labeled by the nucleotides at position *pos* + ℓ .

Consider the set $\mathcal{L}(pos)$ of all labelings of all nodes of tree \mathcal{D} such that the total number of mutations on the edges is minimal. Let m(pos) be the number of labelings in $\mathcal{L}(pos)$ with one mutation between the root and one of its children. Then the contraction cost for position *pos* is defined by: $m(pos)/|\mathcal{L}(pos)|$. Note that, in a parsimonious labeling, there is at most one mutation between the root and its children.

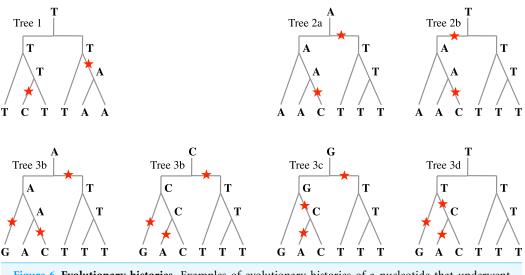


Figure 6 Evolutionary histories. Examples of evolutionary histories of a nucleotide that underwent one duplication followed by two speciation events. In Tree 1, two mutations occurred after speciation: any parsimonious reconstruction from extant nucleotides would agree with this. In Trees 2a and 2b, one mutation occurred before any speciation, and one after: again, any parsimonious reconstruction would agree; however, the ancestral nucleotide could be **A** or **T**. In Trees 3a to 3d, four parsimonious reconstructions with three mutations are given. The number of mutations before speciation can be 0 or 1, and the ancestral nucleotide can be **A**, **C**, **G**, or **T**.

The *contraction cost* of two consecutive segments of length ℓ at position *pos* of the sequence *S* is defined by:

$$C(pos,\ell) = \frac{\sum_{i=pos}^{pos+\ell-1} m(i)/|\mathcal{L}(i)|}{\ell}$$
(2)

We tested this new cost function with the three ortholog tandem repeats ARA, HUM and NEU, in parallel, to obtain the bottom curve of Fig. 5: the positions where the minimal, or near minimal costs are attained are in a small interval, and the contraction costs are nearly halved with respect to the former cost function.

Since the contraction cost for a position, $m(pos)/|\mathcal{L}(pos)|$, is a number between 0 and 1, the total number of mutations as computed by $C(pos, \ell) * \ell$ is not necessarily an integer. When we want to evaluate the number of mutations necessary to perform a contraction, we use the formula:

 $M(pos, \ell) = \lceil C(pos, \ell) * \ell \rceil.$

Contractions and iterations

The labels at the root of different labelings in $\mathcal{L}(pos)$ are the putative ancestors of the nucleotides at position *pos* and *pos* + ℓ , this set of ancestors is denoted by A(pos).

To perform the contraction of two segments of length ℓ on a tandem repeat at position *pos*, we replace the two segments by a single segment. This segment is the sequence of sets

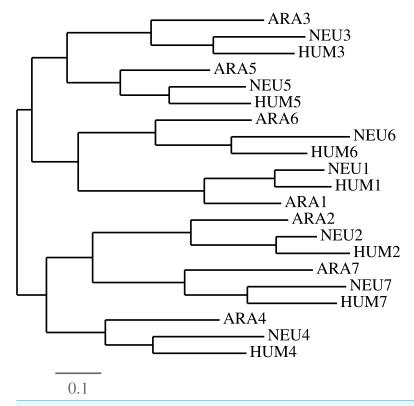


Figure 7 Evidence of duplications followed by speciations. In order to establish the fact that all duplications of the GNB1 beta-propeller preceded the speciation events, we parsed the HUM, NEU and ARA nucleotide sequences in seven 'blades', numbered as in Fig. 4. A phylogenetic tree was then constructed using Dnapars (*Felsenstein, 1993*), showing that the tree supports the fact that duplication events preceded speciation events. Different parsing choices and phylogenetic tools all yielded trees in which the blades grouped by their number rather than by species (data not shown).

A(i) for $pos \le i < pos + \ell$. For example, the contraction of the tandem repeat in the three sequences **TAGTTT**, **CAAATT**, and **TCCATT** used in Fig. 6 would be:

$\{T\}\{A,T\}\{A,C,G,T\}$

After a contraction is performed, it is possible to iterate the procedure by choosing the next more recent duplication. In this case, the leaves of the tree D will be labeled by sets of nucleotides instead of single nucleotides.

Clearly, when using multiple species, iterated contractions may be performed only on *ancient* duplications, that is those which preceded speciation. In the case of GNB1 orthologs, strong evidence supports the fact that all duplications of the WD40 motif preceded all speciations, at least up to the common ancestor of plants and animals. A compelling example of this was shown by Schaper and colleagues in the recent large scale study (*Schaper, Gascuel & Anisimova, 2014*). All approaches are based on the same principle that, when analyzed by phylogenetic tools, the motifs regroup by their position in the tandem sequence rather than by the species they belong to. We provide, in Fig. 7, yet another evidence of this. The contraction procedure may be iterated as often as the sequence is long enough to be contracted. However, it is clear that in some circumstances further contractions are not recommended, as the experiment with the HUM sequence showed in Fig. 5. Before performing a contraction, one should make sure that the interval of positions with minimal costs is small, and that these costs are well below the random costs of performing a contraction anywhere else. The next section contains some further examples of these situations.

Algorithms and complexity

The complete procedure is composed of the following steps, taking as input a motif of known length *p* and a set of *n* ortholog tandem repeats of the motif:

- 1. Create a dataset of *n* tandem repeats of the same length, using insertions and deletions within copies of the motif. If a structure-guided alignment exists, use it.
- 2. Find the values of $(pos, k * p), k \ge 1$ that minimize, or nearly minimize, the contraction cost $C(pos, \ell)$, as defined by Eq. (1) or Eq. (2).
- 3. Decide whether there is sufficient evidence to confirm the position of a further contraction. If not, stop.
- 4. Perform the contraction, and go to Step 2.

Only Steps 2 and 4 can be solved by algorithms, and, fortunately, efficient algorithms are available for these tasks. The computation of contraction costs as defined by Eq. (1) is rather trivial. The computation of $C(pos, \ell)$ as defined by Eq. (2), which calls for the enumeration of all parsimonious labeling of a tree is more complex. Fortunately, there exists an algorithm developed by *Sankoff (1975)* whose data structure can be used to compute both $m(pos)/|\mathcal{L}(pos)|$ and the set of A(pos) of possible ancestors.

Define the cost of a labeling of a tree to be the number of edges that have different labels at their extremities. At each vertex v of the tree \mathcal{D} , we compute, for each symbol b in {**A**, **C**, **T**, **G**, - }, the minimal cost M(v, b) of a labeling of the subtree rooted at v, with base b labeling vertex v; and the number N(v, b) of such minimal cost labelings. A classical dynamic programming algorithm (*Sankoff, 1975*) is used to fill out the structure. The values of these parameters at the root R of \mathcal{D} , and at its left and right children L_R and R_R , are sufficient to compute A(pos) and the contraction cost at position *pos*. namely:

- 1. A(pos) is the set of bases b for which M(R, b) is minimal.
- 2. $|\mathcal{L}(pos)| = \sum_{b \in A(pos)} N(R, b).$
- 3. $m(pos) = |\mathcal{L}(pos)| \sum_{b \in E} N(L_R, b) \cdot N(R_R, b)$, where $E = \{b | b \in A(pos) \text{ and } M(R, b) = M(L_R, b) + M(R_R, b)\}.$

These are all polynomial algorithms, but the extensive computations involved in making tens of thousand of experiments prompted us to implement these with various speed-up techniques (*Lavoie-Mongrain*, 2014).

Contraction	Interval	Width	Costs
1	[383437]	55	0.209-0.219
2	[237305]	69	0.198-0.235
3	[312333]	22	0.229–0.267
4	[58164]	107	0.270-0.306

Table 2 Four contractions. Statistics on the 314 best results from 27,000 four-contractions experiments.

Experimental set-up

We ran $30 \times 30 \times 30 = 27,000$ experiments computing the total cost of the first four contractions yielding the 27,000 data vectors (p_1, p_2, p_3, p_4) and (c_1, c_2, c_3, c_4) , where c_i is the contraction cost at position p_i . The positions p_1 span the 30 lowest costs of a first contraction; for each value of p_1 , the positions p_2 span the 30 lowest costs given a first contraction at p_1 ; for each pair (p_1, p_2) , the positions p_3 span the 30 lowest costs given contractions at p_1 and then p_2 . Position p_4 is the smallest position, for a fourth contraction, with lowest cost. The 72 best results for the first three contractions are available in Table S1. The computation time for the 27,000 experiments was around 21/2 hours (148.5 min) on a 3 GHz processor with 8 GB RAM.

Although we tried contractions of length $k \times 126$, with k > 1, none of them made it in the 30 lowest costs at any iteration.

RESULTS

Using the six ortholog tandem repeats HUM, DAN, DRO, CEL, NEU and ARA, with the phylogeny of Fig. 2, we ran 27,000 experiments minimizing the total cost of going from seven blades to three blades. All scenarios predict four contractions of length 126. This set of experiments allowed us to study the behavior after one, two, three or four contractions.

The 18 optimal scenarios predicted a total of 122 mutations to perform the four contractions, and 296 near optimal scenarios predicted a total of 123 mutations. Table 3 sums up the principal characteristics of these 314 scenarios giving, for each contraction, the interval of positions in which minimal or near minimal values were obtained, its width, and the range of its costs.

Since the cost of the fourth contraction is in the upper range of cost values (an example using one of the optimal scenarios is shown in Fig. 8), and is predicted at positions ranging from 58 to 164, which is a very wide interval given that the period is 126, we halted the reconstruction process before the fourth contraction, leading us to concentrate on scenarios with three contractions yielding a 4-bladed ancestor.

The 4-bladed ancestor

Of the 27,000 experiments, 6 optimal scenarios predicted a total of 83 mutations to perform three contractions, and 66 near optimal scenarios predicted a total of 84 mutations. Table 3 shows the same statistics as Table 2 for these 72 scenarios. The complete results for these experiments are in Table S1.

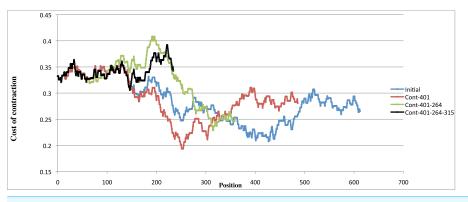


Figure 8 Variation of costs over four contractions. Example of the variation of the cost of four contractions using one scenario with minimal cost. The blue curve shows, for each position of the input sequence, the cost of performing a contraction of length 126 at that position. The red curve shows the contraction costs computed on the sequence obtained by contracting the input sequence at position 401. The green curve shows the contraction costs computed on the sequence obtained by contracting the input sequence at position 401 then at position 264. The black curve shows the contraction costs computed on the sequence obtained by contracting the input sequence at position 315. While the first three curves have well defined minima, the fourth curve is in the upper range of all curves, showing that performing a fourth contraction might be meaningless.

Contraction	Interval	Width	Costs
1	[383402]	20	0.209-0.218
2	[251264]	14	0.188-0.221
3	[312322]	11	0.228-0.253

Table 3 Three contractions. Statistics on the 72 best results from 27,000 three-contractions experiments.

We analyzed all 72 reconstructions of the 4-bladed ancestor corresponding to the 72 best results. These reconstructions are remarkably similar, which is not completely surprising given the narrow ranges of possible first, second and third contraction, as can be seen in Table 3, and Fig. S1.

For each triplet of positions, we performed the three contractions, and then translated the contracted part using the extended one-letter amino acid code (*Cornish-Bowden*, 1985) that assigns the letter **B** to the sets of amino acids {**D**, **N**}, and the letter **Z** to the set {**E**, **Q**}, see Fig. 9.

All reconstructions left undisturbed both the first two repeats—positions 1–251—and the last 39 positions (see Fig. S1). All reconstructions agreed on the following translated middle part sequence (see Fig. S2), where dots indicate unresolved amino acids:

GHTG.V....F.PBD.....G..D.TC..WD...GZ.V....H.N.....FS..G

and 69 sequences (96%) agreed with the pattern:

GHTG.V....F.PBD..L..G..D.TC..WD...GZ.V.....H.N......FS..G

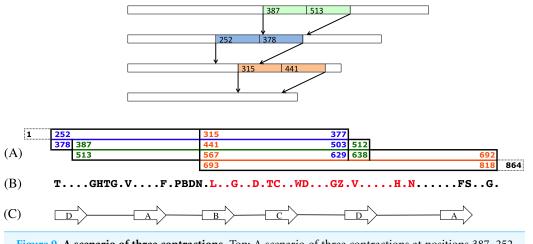


Figure 9 A scenario of three contractions. Top: A scenario of three contractions at positions 387, 252 and 315. (A) This diagram shows the positions of the original sequence that have been contracted. The first one, in green, is the contraction of the intervals [387..512] and [513..638]. The second one, in blue and overlapping the green contraction, is the contraction of intervals [252..377] and [378..503]. The last contraction, in orange and overlapping the two other contractions, starts at position 315, and runs up to the original position 818, leaving 46 non contracted nucleotides, in dashed lines, at the end of the sequences. (B) This is the reconstructed amino acids sequence of the contracted segment. Amino acids **B** and **Z** stand, respectively, for the sets of amino acids {**D**, **N**} and {**E**, **Q**} (*Cornish-Bowden, 1985*). The amino acids in red correspond to the pattern that was found in the Slime mold GNB1 protein (see text). (C) The positions of the strands A, B, C and D in the contracted sequence.

Searching these patterns in various protein databases yields no result. However, searching for the longest (sub)pattern that has at least one hit, gives very interesting results. This pattern is:

L..G..D.TC..WD...GZ.V....H.N

The pattern was not found in the six sequences that were chosen for the current experiment. However, there were 32 hits in the TrEMBL protein database (*Bairoch & Apweiler, 2000*) using ScanProsite (*de Castro et al., 2006*). Of these hits, two were for the slime molds *Dictyostelium discoideum* and *Dictyostelium purpureum*, whose divergence is presumed to date back to 400 million years (*Parikh et al., 2010*). The rest of the hits are all Basidiomycota, which is a fungus phylum: these hits include mostly rots and smuts, with some yeasts, and two edible mushrooms, the *Shitake* and the *Champignon de Paris*.

DISCUSSION AND CONCLUSIONS

While phylogenetic tools and data have been used for tandem repeat reconstruction, it was always in the context of the fixed boundaries hypothesis. The procedure that we presented in this paper constitutes, to the best of our knowledge, the first attempt to systematically use phylogenetic data to reconstruct the evolutionary history of tandem repeats with dynamic boundaries. The choice of parsimony was based on two considerations. First, we are generalizing a parsimony-based algorithm, and, second, since we are using a very computationally intensive algorithm, efficiency is an issue: computation of the cost functions is at the core of the algorithm, it has to be simple. In the future, it would be interesting to define more sophisticated cost functions since we now know the value of basic parameters, such as the period.

In a recent paper (*Schaper, Gascuel & Anisimova, 2014*), the authors found 3,091 tandem repeats in human proteins, of these, 17% are conserved since the root of the vertebrates. This means that there are potentially 525 sequences that could be studied by techniques presented in this paper. The phylogenetic methods proposed in that paper, based on a fixed boundary tandem repeats, are much more efficient than our exhaustive search approach and should be used whenever the fixed boundary hypothesis is suspected to apply, such as when the repeated units are bound by exons (*Schaper, Gascuel & Anisimova, 2014*). However, as the example of WD40 beta-propeller tandem repeat suggests, duplication histories with dynamic boundaries should not be readily excluded.

Are the proposed 4-bladed ancestor, or the common motif that we identified, realistic? It is hard to answer yes or no, but a few good arguments plead in favor of this ancestor. Taking an *a posteriori* look back at Fig. 5, in which we studied the species separately, one can see that the NEU sequence is a very good predictor for the consensus position of the most recent duplication. It is also the most simple species in our dataset, and the common motif that we found points to early branching or slow evolving organisms, such as slime molds.

Given the fact that the first repeat interacts with the last one in order to close the circular structure of the beta-propeller, as seen in Fig. 1, it seems very plausible that the expansion of the number of blades in propellers should leave the beginning and end of the sequence undisturbed, since they probably need co-evolution to properly glue the structure. This may explain why we could not reliably predict any contraction overlapping the first two repeats, or the last 12 residues.

Most importantly, no overwhelming evidence supports the hypothesis of fixed boundaries in the WD40 beta-propeller tandem repeat. As Fig. 9C shows, while the boundaries of the second most recent duplication coincide with the boundaries of a *DABC* parsing of the WD40 motif, the boundaries of the most recent are a little bit offset from those of the *DABC* parsing, and the boundaries of the third most recent do not agree at all with those of the *DABC* parsing. This pattern is true for all the 72 best results.

ADDITIONAL INFORMATION AND DECLARATIONS

Funding

Anne Bergeron is partially supported by NSERC Grant number 05729-2014. The funders had no role in study design, data collection and analysis, decision to publish, or preparation of the manuscript.

Grant Disclosures

The following grant information was disclosed by the authors: NSERC: 05729-2014.

Competing Interests

Anne Bergeron is an Academic Editor for PeerJ.

Author Contributions

- Philippe Lavoie-Mongrain performed the experiments, performed the computation work, reviewed drafts of the paper.
- Mahdi Belcaid conceived and designed the experiments, performed the experiments, analyzed the data, performed the computation work, reviewed drafts of the paper.
- Aïda Ouangraoua conceived and designed the experiments, performed the experiments, analyzed the data, prepared figures and/or tables, performed the computation work, reviewed drafts of the paper.
- Anne Bergeron conceived and designed the experiments, analyzed the data, wrote the paper, prepared figures and/or tables, reviewed drafts of the paper.
- Guylaine Poisson conceived and designed the experiments, analyzed the data, prepared figures and/or tables, reviewed drafts of the paper.

Supplemental Information

Supplemental information for this article can be found online at http://dx.doi.org/ 10.7717/peerj-cs.6#supplemental-information.

REFERENCES

- Altschul SF, Gish W, Miller W, Myers EW, Lipman DJ. 1990. Basic local alignment search tool. Journal of Molecular Biology 215:403–410 DOI 10.1016/S0022-2836(05)80360-2.
- Andrade MA, Perez-Iratxeta C, Ponting CP. 2001. Protein repeats: structures, functions, and evolution. *Journal of Structural Biology* **134**(2–3):117–131 DOI 10.1006/jsbi.2001.4392.
- Bairoch A, Apweiler R. 2000. The swiss-prot protein sequence data bank and its supplement trembl in 2000. *Nucleic Acids Research* 28:45–48 DOI 10.1093/nar/28.1.45.
- **Benson G, Dong L. 1999.** Reconstructing the duplication history of a tandem repeat. In: *Proceedings of the seventh international conference on intelligent systems for molecular biology*. Palo Alto: AAAI Press, 44–53.
- Chaudhuri I, Soding J, Lupas AN. 2008. Evolution of the beta-propeller fold. *Proteins* 71(2):795–803 DOI 10.1002/prot.21764.
- Cornish-Bowden A. 1985. Nomenclature for incompletely specified bases in nucleic acid sequences: recommendations 1984. *Nucleic Acids Research* 13(9):3021–3030 DOI 10.1093/nar/13.9.3021.
- de Castro E, Sigrist CJ, Gattiker A, Bulliard V, Langendijk-Genevaux PS, Gasteiger E, Bairoch A, Hulo N. 2006. ScanProsite: detection of PROSITE signature matches and ProRule-associated functional and structural residues in proteins. *Nucleic Acids Research* 34(Suppl 2):W362–W365 DOI 10.1093/nar/gkl124.
- Dunn CW, Hejnol A, Matus DQ, Pang K, Browne WE, Smith SA, Seaver E, Rouse GW, Obst M, Edgecombe GD, Sørensen MV, Haddock SH, Schmidt-Rhaesa A, Okusu A, Kristensen RM, Wheeler WC, Martindale MQ, Giribet G. 2008. Broad phylogenomic sampling improves resolution of the animal tree of life. *Nature* 452(7188):745–748 DOI 10.1038/nature06614.

- Edgar RC. 2004. MUSCLE: multiple sequence alignment with high accuracy and high throughput. *Nucleic Acids Research* 32(5):1792–1797 DOI 10.1093/nar/gkh340.
- Elias I. 2006. Settling the intractability of multiple alignment. *Journal of Computational Biology* 13(7):1323–1339 DOI 10.1089/cmb.2006.13.1323.
- Felsenstein J. 1993. Phylip (phylogeny inference package) version 3.5c. Distributed by the author. Department of Genetics, University of Washington, Seattle. Available at http://cmgm.stanford. edu/phylip/#1.
- Flicek P, Amode M, Barrell D. 2014. Ensembl 2014. *Nucleic Acids Research* 42(D1):D749–D755 DOI 10.1093/nar/gkt1196.
- **Gascuel O, Bertrand D, Elemento O. 2005.** Reconstructing the duplication history of tandemly repeated sequences. In: Gascuel O, ed. *Mathematics of evolution and phylogeny*. Oxford: Oxford University Press, 205–235.
- Hanson RM, Prilusky J, Renjian Z, Nakane T, Sussman JL. 2013. Jsmol and the next-generation web-based representation of 3D molecular structure as applied to proteopedia. *Israel Journal of Chemistry* 53(3–4):207–216 DOI 10.1002/ijch.201300024.
- Jaitly D, Kearney P, Lin G, Ma B. 2002. Methods for reconstructing the history of tandem repeats and their application to the human genome. *Journal of Computer and System Sciences* 65(3):494–507 DOI 10.1016/S0022-0000(02)00006-5.
- Kajava AV. 2012. Tandem repeats in proteins: from sequence to structure. *Journal of Structural Biology* 179(3):279–288 DOI 10.1016/j.jsb.2011.08.009.
- Lavoie-Mongrain P. 2014. Reconstruction des ancêtres de séquences partageant leur histoire duplicative. Master's Thesis, Université du Québec à Montréal. *Available at http://www.archipel.uqam.ca/*.
- Li D, Roberts R. 2001. WD-repeat proteins: structure characteristics, biological function, and their involvement in human diseases. *Cellular and Molecular Life Science* **58**(14):2085–2097 DOI 10.1007/PL00000838.
- Matroud AA, Tuffley CP, Bryant D, Hendy MD. 2012. A comparison of three heuristic methods for solving the parsing problem for tandem repeats. In: *Advances in bioinformatics and computational biology*. Berlin: Springer, 37–48.
- Parikh A, Miranda ER, Katoh-Kurasawa M, Fuller D, Rot G, Zagar L, Curk T, Sucgang R, Chen R, Zupan B, Loomis WF, Kuspa A, Shaulsky G. 2010. Conserved developmental transcriptomes in evolutionarily divergent species. *Genome Biology* 11(3):R35 DOI 10.1186/gb-2010-11-3-r35.
- **Rivals E. 2004.** A survey on algorithmic aspects of tandem repeats evolution. *International Journal of Foundations of Computer Science* **15(2)**:225–257 DOI 10.1142/S012905410400239X.
- Sammeth M, Stoye J. 2006. Comparing tandem repeats with duplications and excisions of variable degree. *IEEE/ACM Transactions on Computational Biology and Bioinformatics* 3(4):395–407 DOI 10.1109/TCBB.2006.46.
- Sankoff D. 1975. Minimal mutation trees of sequences. *SIAM Journal on Applied Mathematics* 28:35–42 DOI 10.1137/0128004.
- Schaper E, Gascuel O, Anisimova M. 2014. Deep conservation of human protein tandem repeats within the eukaryotes. *Molecular Biology and Evolution* **31**(5):1132–1148 DOI 10.1093/molbev/msu062.
- Sievers F, Wilm A, Dineen D, Gibson TJ, Karplus K, Li W, Lopez R, McWilliam H, Remmert M, Soding J, Thompson JD, Higgins DG. 2011. Fast, scalable generation of high-quality

protein multiple sequence alignments using Clustal Omega. *Molecular Systems Biology* 7:539 DOI 10.1038/msb.2011.75.

- Stirnimann CU, Petsalaki E, Russell RB, Muller CW. 2010. WD40 proteins propel cellular networks. *Trends in Biochemical Sciences* 35(10):565–574 DOI 10.1016/j.tibs.2010.04.003.
- Tremblay Savard O, Bertrand D, El-Mabrouk N. 2011. Evolution of orthologous tandemly arrayed gene clusters. *BMC Bioinformatics* 12(Suppl 9):S2 DOI 10.1186/1471-2105-12-S9-S2.
- Zhuo D, Madden R, Elela SA, Chabot B. 2007. Modern origin of numerous alternatively spliced human introns from tandem arrays. *Proceedings of the National Academy of Sciences of the United States of America* 104(3):882–886 DOI 10.1073/pnas.0604777104.